

Image interdisciplinaire de la Combinatoire Analytique

Interdisciplinary image of Analytic Combinatorics

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Sous direction de
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Why analytic combinatorics

Outline of the current talk

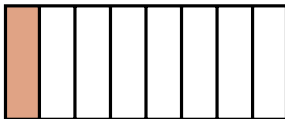
Interdisciplinary image of Analytic Combinatorics

Part I

Part II

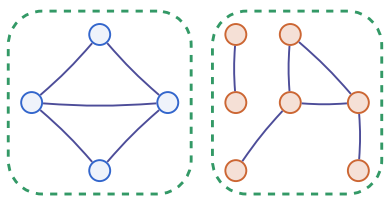
Part I. Analytic Combinatorics

Level I. Symbolic method



(You can never be too thorough)

Known convolution rules



convolution

generating function

name

$$\sum_{k=0}^n a_k b_{n-k}$$

$$A(z) := \sum_{n \geq 0} a_n z^n$$

ordinary GF

$$\sum_{k=0}^n \binom{n}{k} a_k b_{n-k}$$

$$A(z) := \sum_{n \geq 0} a_n \frac{z^n}{n!}$$

exponential GF

$$\sum_{d|n} a_d b_{n/d}$$

$$A(s) := \sum_{n \geq 1} \frac{a_n}{n^s}$$

Dirichlet series

$$\sum_{k=0}^n \binom{n}{k} a_k b_{n-k} 2^{k(n-k)}$$

$$A(z) := \sum_{n \geq 0} a_n \frac{z^n}{2^{\binom{n}{2}} n!}$$

graphic GF

Combinatorial operations

Operation	interpretation
$A(z) + B(z)$	disjoint union
$A(z)B(z)$	cartesian product
$\exp(A(z))$	set
$A(B(z))$	substitution
$z\partial_z A(z)$	pointing
$A(z) \odot B(z)$	Hadamard product
...	...

The philosophy of the symbolic method. (Bergeron, Labelle, Leroux).

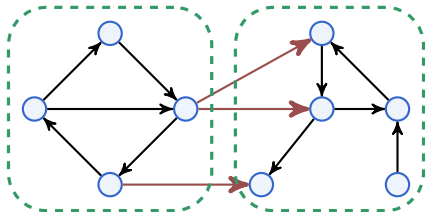
Decomposition \Rightarrow Equation

Follow-up: asymptotic analysis. (Flajolet, Odlyzko, ...).

Equation \Rightarrow GF expansion \Rightarrow Asymptotics

Graphic convolution rule

$$A(z) = a_0 + \frac{a_1 z}{1! 2^{\binom{1}{2}}} + \frac{a_2 z^2}{2! 2^{\binom{2}{2}}} + \dots, \quad B(z) = b_0 + \frac{b_1 z}{1! 2^{\binom{1}{2}}} + \frac{b_2 z^2}{2! 2^{\binom{2}{2}}} + \dots$$



coefficient level

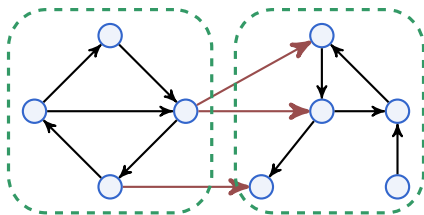
graphic GF level

$$c_n := \sum_{k=0}^n \binom{n}{k} a_k b_{n-k} 2^{k(n-k)}$$

$$C(z) = A(z) \cdot B(z)$$

Graphic convolution rule

$$A(z, \mathbf{w}) = a_0 + \frac{a_1(\mathbf{w})z}{1!(1+\mathbf{w})^{\binom{1}{2}}} + \frac{a_2(\mathbf{w})z^2}{2!(1+\mathbf{w})^{\binom{2}{2}}} + \frac{a_3(\mathbf{w})z^3}{3!(1+\mathbf{w})^{\binom{3}{2}}} + \dots$$



coefficient level

graphic GF level

$$c_n := \sum_{k=0}^n \binom{n}{k} a_k b_{n-k} 2^{k(n-k)}$$

$$C(z) = A(z) \cdot B(z)$$

$$c_n := \sum_{k=0}^n \binom{n}{k} a_k(\mathbf{w}) b_{n-k}(\mathbf{w}) (1+\mathbf{w})^{k(n-k)}$$

$$C(z, \mathbf{w}) = A(z, \mathbf{w}) B(z, \mathbf{w})$$

\mathbf{w} marks edges

Enumerative result

[De Panafieu, D. '19]

Exponential Hadamard product:

$$\left(\sum_{n \geq 0} a_n \frac{z^n}{n!} \right) \odot \left(\sum_{n \geq 0} b_n \frac{z^n}{n!} \right) := \sum_{n \geq 0} a_n b_n \frac{z^n}{n!}$$

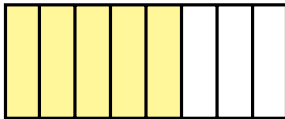
Exponential GF for graphs, graphic GF for sets:

$$\mathbb{G}(z, w) = \sum_{n \geq 0} (1+w)^{\binom{n}{2}} \frac{z^n}{n!}, \quad \widehat{\text{Set}}(z, w) = \sum_{n \geq 0} \frac{1}{(1+w)^{\binom{n}{2}}} \frac{z^n}{n!},$$

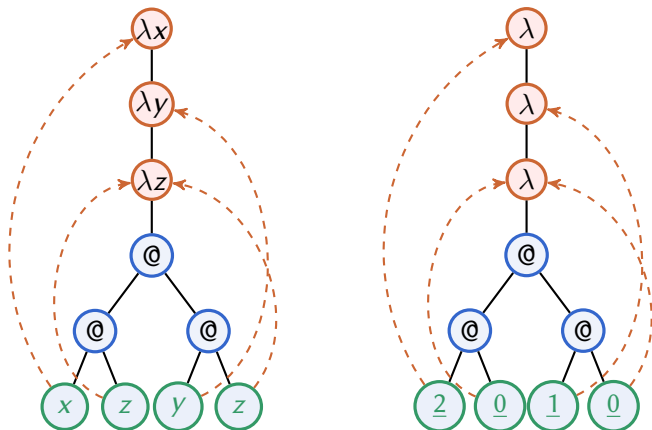
Theorem. (Rediscovery of results by Liskovets and Robinson)

	digraphs	graphs
EGF, connected	$-\log \left(\mathbb{G}(z, w) \odot \frac{1}{\mathbb{G}(z, w)} \right)$	$-\log \left(\frac{1}{\mathbb{G}(z, w)} \right)$
components from given family \mathcal{C}	$GGF = \frac{1}{\widehat{\text{Set}}(z, w) \odot e^{-C(z, w)}}$	$EGF = \frac{1}{e^{-C(z, w)}}$

Level II. Systems of algebraic equations



Lambda terms



A tree-like representation of a *closed* lambda term
 $\lambda x. \lambda y. \lambda z. (x z)(y z)$

Drmotá–Lalley–Woods theorem (simplified)

Theorem.

- ▶ Let $F(z)$ be a generating function satisfying

$$F(z) = \Phi(F(z), z)$$

with Φ is nonlinear, strongly-connected and has “combinatorial origin”

- ▶ Then, there exists $\rho > 0$ such that as $z \rightarrow \rho$, $z \in \mathbb{C}$

$$F(z) \sim a - b\sqrt{1 - \frac{z}{\rho}}$$

Example. Plain lambda terms:



$$L(z) = zL(z) + zL^2(z) + \frac{z}{1-z}$$

Drmotá–Lalley–Woods theorem in infinite dimension

$$F(z, \vec{u}) = \Phi(F(z, \vec{u}), z, \vec{u})$$

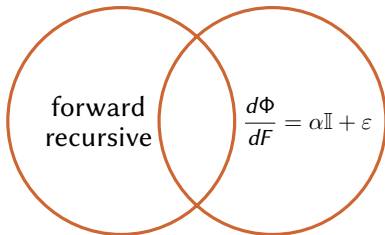
Theorem. [Drmotá, Gittenberger, Morgenbesser '2012] simplified

Let $\dim \Phi = \infty$. If the Jacobian $\frac{d\Phi}{dF}$ can be represented as $\alpha \cdot \mathbb{I} + \varepsilon$, where ε is a compact operator, then DLW expansion holds:

$$F(z) \sim a(\vec{u}) - b(\vec{u}) \sqrt{1 - \frac{z}{\rho(\vec{u})}}$$

Theorem. [Bendkowski, Bodini, D. '2019] simplified

If Φ is ∞ -dimensional *forward-recursive* then DLW expansion holds.



Formulation of our result

[Bendkowski, Bodini, D.]

Definition. (Forward recursive systems)

$$\vec{L}_m = \vec{K}_m(\vec{L}_m, \vec{L}_{m+1}, z, \vec{u}), \quad m = 0, 1, 2, \dots, \infty.$$

satisfies the conditions

- ▶ $\vec{K}_m(\vec{\ell}_1, \vec{\ell}_2, z, \vec{u}) \succeq \vec{K}_\infty(\vec{\ell}_1, \vec{\ell}_2, z, \vec{u})$
- ▶ \vec{K}_∞ is DLW-well-defined
- ▶ $\vec{K}_\infty(\vec{L}_\infty, \vec{L}_\infty, z, \vec{u}) - \vec{K}_m(\vec{L}_\infty, \vec{L}_\infty, z, \vec{u}) \preceq \vec{A}(z, \vec{u}) \cdot \underbrace{B(z, \vec{u})^m}_{<1 \text{ at } z=\rho, \vec{u}=\vec{1}}$;

Theorem. (Coefficient transfer)

$$[z^n] \vec{L}_m(z, \vec{u}) \sim [z^n] \sum_{k \geq 0} \vec{c}_{m,k} \left(1 - \frac{z}{\rho(\vec{u})}\right)^{k/2}$$

$$\|\vec{c}_{m,k} - \vec{c}_{\infty,k}\| \sim C^{-m}$$

Illustration of a forward-recursive system

Example. *Closed lambda terms.*

$$L_0(z) = zL_1(z) + zL_0(z)^2 ,$$

$$L_1(z) = zL_2(z) + zL_1(z)^2 + z ,$$

$$L_2(z) = zL_3(z) + zL_2(z)^2 + z + z^2 ,$$

...

$$L_\infty(z) = zL_\infty(z) + zL_\infty(z)^2 + \frac{z}{1-z} .$$

Conditions to check.

- ▶ $z\lambda_1 + z\lambda_2^2 + z + z^2 + \dots + z^m \preceq z\lambda_1 + z\lambda_2^2 + \frac{z}{1-z}$
- ▶ $z^{m+1} + \dots \preceq \alpha(z) \cdot \beta(z)^m$
- ▶ $L_\infty(z)$ satisfies DLW conditions

Generalisation. Decorate with marking variables.

Multiparametric sampling

Context-free grammar with terminals $\bullet_1, \bullet_2, \bullet_3, \bullet_4$ and non-terminals S_1, \dots, S_n :

$$S_i \rightarrow \sum_j T_{ij}(S_1, \dots, S_n, \bullet_1, \bullet_2, \bullet_3, \bullet_4)$$

Exact sampling. Given positive integers (n_1, n_2, \dots, n_d) , sample *uniformly at random* a word $w \in S_1$ such that

$$|w|_{\bullet_1} = n_1, \quad |w|_{\bullet_2} = n_2, \quad \dots, \quad |w|_{\bullet_d} = n_d$$

:(This problem is #P-complete):

Boltzmann sampling. Given positive reals (z_1, z_2, \dots, z_d) , sample a word $w \in S_1$ so that

$$\mathbb{P}(w : |w|_{\bullet_1} = n_1, |w|_{\bullet_2} = n_2, \dots, |w|_{\bullet_d} = n_d) = \frac{z_1^{n_1} z_2^{n_2} \dots z_d^{n_d}}{\sum_{n_1, \dots, n_d} z_1^{n_1} z_2^{n_2} \dots z_d^{n_d}}$$

(: There exists a polynomial algorithm :)

Boltzmann sampler

$$S_i \rightarrow \sum_j T_{ij}(S_1, \dots, S_n, \bullet)$$

Algorithm 1: Boltzmann sampler for context-free grammars

Data: real value $z > 0$

Result: Random word from Boltzmann distribution

Function $\Gamma R_i(z)$:

if R_i is terminal **then**

return \bullet ;

for all j **do**

$p_j := \frac{T_{ij}(S_1(z), \dots, S_n(z), z)}{R_i(z)}$;

 Choose the transition T_{ij} with probability p_j ;

$A_1 A_2 \dots A_k := T_{ij}$;

return $\Gamma A_1(z) \Gamma A_2(z) \dots \Gamma A_k(z)$;

Multiparametric Boltzmann sampler

$$S_i \rightarrow \sum_j T_{ij}(S_1, \dots, S_n, \bullet_1, \bullet_2, \dots, \bullet_\ell)$$

Algorithm 2: Boltzmann sampler for context-free grammars

Data: real values $z_1, z_2, \dots, z_\ell > 0$

Result: Random word from Boltzmann distribution

Function $\Gamma R_i(z)$:

if R_i is terminal \bullet_k **then**

return \bullet_k ;

for all j **do**

$p_j := \frac{T_{ij}(S_1(z), \dots, S_n(z), z_1, z_2, \dots, z_\ell)}{R_i(z)}$;

 Choose the transition T_{ij} with probability p_j ;

$A_1 A_2 \dots A_k := T_{ij}$;

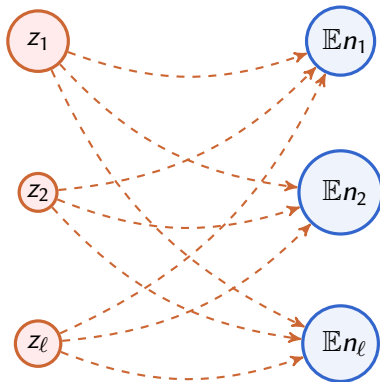
return $\Gamma A_1(z) \Gamma A_2(z) \dots \Gamma A_k(z)$;

Tuning of the multiparametric Boltzmann sampler

Handles \mathbf{z}

\Rightarrow

Expectations $\mathbb{E}n_k$



!! The handles cannot be tuned independently !!

Tuning equation:

$$z_1 \frac{\partial_{z_1} F(\vec{z})}{F(\vec{z})} = \mathbb{E}n_1, \quad z_2 \frac{\partial_{z_2} F(\vec{z})}{F(\vec{z})} = \mathbb{E}n_2, \quad \dots, \quad z_\ell \frac{\partial_{z_\ell} F(\vec{z})}{F(\vec{z})} = \mathbb{E}n_\ell$$

Tuning of the multiparametric Boltzmann sampler

[Bendkowski, Bodini, D.]

Theorem. Let the expected values n_1, \dots, n_ℓ of the terminals

$\bullet_1, \bullet_2, \dots, \bullet_\ell$ be given. Let $S_k(\mathbf{z})$ satisfy

$$S_1 = \Phi_1(S_1, \dots, S_n, \mathbf{z}),$$

...

$$S_n = \Phi_n(S_1, \dots, S_n, \mathbf{z}).$$

The tuning vector $(z_1, \dots, z_\ell) = (e^{x_1}, \dots, e^{x_\ell})$ can be obtained by solving a convex optimisation problem

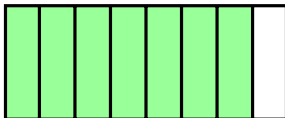
$$S_1 - n_1 x_1 - n_2 x_2 - \dots - n_\ell x_\ell \rightarrow \min_{(S_1, \dots, S_n, x_1, \dots, x_\ell)},$$

$$S_1 \geq \log \Phi_1(e^{S_1}, \dots, e^{S_n}, e^{x_1}, \dots, e^{x_\ell}),$$

...

$$S_n \geq \log \Phi_n(e^{S_1}, \dots, e^{S_n}, e^{x_1}, \dots, e^{x_\ell}).$$

Level III. Airy function and saddle point



Boltzmann tuning as saddle point equation

$$[z^n]f(z) = \frac{1}{2\pi i} \oint_{|z|=R} e^{g(z)} \frac{dz}{z}$$

\Downarrow

$$g'(z) = 0 \quad \Leftrightarrow \quad z \frac{f'(z)}{f(z)} = n$$

\Downarrow

$$\log f(e^z) - nz \rightarrow \min_z \quad \Leftrightarrow \quad \mathbb{E}_z[N] = n$$

$$\left(\mathbb{P}_x(N = k) = \frac{x^k [z^k] f(z)}{f(x)} \right)$$

Coalescence of saddle points

Gaussian case.

$$\begin{aligned} [z^n]f(z) &= \frac{1}{2\pi i} \oint_{|z|=R} e^{g(z)} \frac{dz}{z} \approx \frac{1}{2\pi i} \oint_{|z|=R} e^{g(z_0) + \frac{1}{2}g''(z_0)(z-z_0)} \frac{dz}{z} \\ &\approx \frac{1}{\sqrt{2\pi g''(z_0)}} e^{g(z_0)} \end{aligned}$$

Coalescence of saddle points. When $g'(z) = g''(z) = 0$,

$$g(z) \approx g(z_0) + \frac{1}{6}g'''(z_0)(z - z_0)^3$$

Airy function.

$$A(y) = \int_{\mathbb{R}} e^{i(x^3/3 - xy)} dx, \quad A''(y) + yA(y) = 0$$

FJKLP lemma for graphs

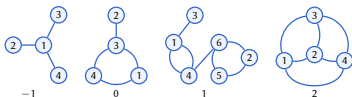
[Flajolet, Janson, Knuth, Łuczak, Pittel] → Airy function is hidden inside

Point of saddle coalescence. Graphs with n vertices, m edges,
 $m = \frac{n}{2}(1 - \mu n^{-1/3}), \mu \rightarrow \infty$:

$$\frac{n!}{|\mathcal{G}_{n,m,\Delta}|} \frac{1}{2\pi i} \oint \frac{U(z)^{n-m}}{(n-m)!} e^{\underset{\substack{\uparrow \\ \text{unicycles}}}{V(z)}} \frac{dz}{z^{n+1}} = 1 - \frac{5}{24|\mu|^3} + O(\mu^{-6})$$

\uparrow graphs \uparrow consist of \uparrow trees

Refined analysis.



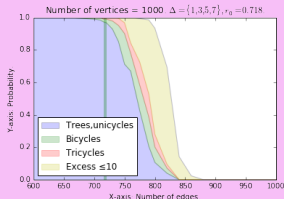
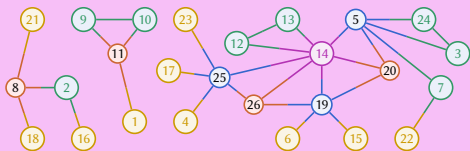
$$\frac{n!}{|\mathcal{G}_{n,m,\Delta}|} \frac{1}{2\pi i} \oint \frac{U(z)^{n-m}}{(n-m)!} e^{\underset{\substack{\uparrow \\ \text{unicycles}}}{V(z)}} \frac{c}{(1 - T(z))^{3\gamma}} \frac{dz}{z^{n+1}} \sim \frac{c}{|\mu|^{3\gamma}}$$

\uparrow graphs \uparrow consist of \uparrow trees \uparrow complex comp.

Graphs with degree constraints

[De Panafieu, Ramos], [D., Ravelomanana]

Allowed degrees: $\Delta = \{\delta_1, \delta_2, \delta_3, \delta_4, \delta_5, \dots\}$, $\delta_1 = 1$ or $\delta_2 = 1$.



EGF of allowed degrees helps to construct the trees and unicycles, next higher excess

$$\omega(z) = \sum_{d \in \Delta} \frac{z^d}{d!} \Rightarrow T(z) = z\omega'(T(z)), \quad U(z) = z\omega(T) - \frac{T(z)^2}{2}$$

Point of saddle coalescence. Graphs with n vertices, m edges, $m = \alpha n(1 - \mu n^{-1/3})$, $\mu \rightarrow \infty$,

$$\widehat{z} \frac{\omega''(\widehat{z})}{\omega'(\widehat{z})} = 1, \quad \widehat{z} \frac{\omega'(\widehat{z})}{\omega(\widehat{z})} = 2\alpha.$$

2-SAT contradictory component

[Kim '2008], [D. '2019+]

Point of saddle coalescence. 2-CNF with n variables, m clauses,
 $m = n(1 - \mu n^{-1/3})$, $\mu \rightarrow \infty$:

$$\mathbb{P}(F_{n,m} \text{ is satisfiable}) \sim 1 - \frac{1}{16|\mu|^3}$$

Theorem. [D. '2019] Cubic contradictory cores of excess r appear with probability $\Theta(|\mu|^{-3r})$.

Open question. [De Panafieu, D., Ravelomanana '2019]

Digraphs with n vertices, m oriented edges, $m = n(1 - \mu n^{-1/3})$,
 $\mu \rightarrow \infty$:

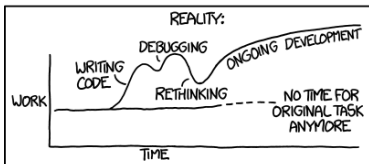
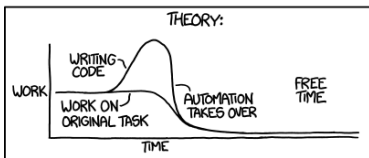
$$\mathbb{P}\left(\begin{array}{l} \text{every strong component is} \\ \text{an isolated vertex or a cycle} \end{array}\right) \sim 1 - \frac{C}{|\mu|^3}, \quad C = ?$$

Part II. Interdisciplinary image

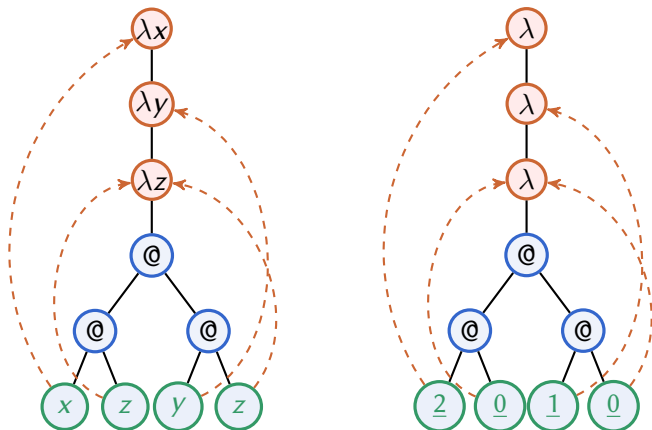
Level I. Software verification



"I SPEND A LOT OF TIME ON THIS TASK.
I SHOULD WRITE A PROGRAM AUTOMATING IT!"



Lambda terms



A tree-like representation of a *closed* lambda term
 $\lambda x. \lambda y. \lambda z. (x z) (y z)$

Lambda terms

- ▶ Church–Turing thesis

Lambda terms \Leftrightarrow Turing machines

Beta-reduction step \Leftrightarrow Crank turn

- ▶ Curry–Howard correspondence

Typed lambda terms \Leftrightarrow Proofs in Hilbert-style deduction logic

Types \Leftrightarrow Statements

THESIS FOR THE DEGREE OF LICENTIATE OF PHILOSOPHY

Testing an Optimising Compiler by Generating Random Lambda Terms

MICHAŁ H. PAŁKA

CHALMERS | GÖTEBORG UNIVERSITY



Department of Computer Science and Engineering
CHALMERS UNIVERSITY OF TECHNOLOGY AND GÖTEBORG UNIVERSITY
Göteborg, Sweden 2012

- ▶ CSmith ['2011]: generation of random C programs, 325 previously unreported bugs in GCC, LLVM, other C compilers and CompCert (with a formally verified core)
- ▶ Michał Pałka et al. ['2012]: Eight failures and four bugs detected, reported and fixed (GHC optimising compiler)

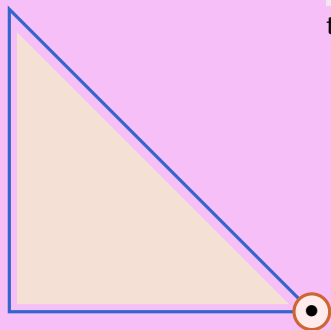
Statistical properties of closed lambda terms

[Bendkowski, Bodini, D. '2019]

Parameter	Mean, \sim		Distribution	
	plain	closed	plain	closed
Variables	0.307 <i>n</i>		Normal	
Abstractions	0.258 <i>n</i>		Normal	
Successors	0.129 <i>n</i>		Normal	
Redexes	0.091 <i>n</i>		Normal	
Index value	0.420		Geometric	
Redex search time	6.222	6.054	Discrete	Discrete
Head abstractions	0.420	1.447	Geometric	Discrete
<i>m</i> -openness	2.019	0	Discrete	trivial
Free variables	5.722	0	Discrete	trivial
Unary height profile	0.122 \sqrt{n}		Rayleigh	
Natural height profile	0.412 \sqrt{n}		Rayleigh	

Random generation of closed lambda terms

[Bendkowski, Bordini, D. '2019]



Theorem. Polynomial algorithm for tuning with given

- ▶ number of atomic nodes of distinguished colors
- ▶ number of **redexes** (i.e. patterns necessary to perform a computation step in lambda calculus)
- ▶ number of head abstractions
- ▶ number of closed subterms
- ▶ number of any tree-like patterns

Having *corner cases* instead of *uniform distribution* is necessary.

Multiparametric Tuner Implementation

[Bendkowski, Bodini, D.]



Docs » Paganini's documentation [Edit on GitHub](#)

Paganini's documentation

Paganini is a lightweight python library for tuning multiparametric combinatorial specifications.

Given a combinatorial specification, expressed using a domain-specific language closely resembling Flajolet and Sedgewick's *symbolic method*, Paganini gives its users some additional control over the distribution of structures constructed using the designed samplers.

```
>>> from paganini import *
>>> spec = Specification()
>>> z, u, M = Variable(1000), Variable(200), Variable()
>>> spec.add(M, z + u * z * M + z * M ** 2)
>>> spec.run_tuner(M)
```

Level II. Data sciences



Definition. A *network* is a graph or directed graph, formed by *real-world data*.

Differences exhibited (networks versus Erdős–Rényi graphs)

- ▶ Density
- ▶ Degree distribution
- ▶ Diameter, shortest path
- ▶ Connectivity, clustering
- ▶ Spectral properties
- ▶ other metrics

Degree constraints approach

[De Panafieu, Ramos], [D., Ravelomanana]

Set of degree constraints

$$\Delta = \{\delta_1, \delta_2, \delta_3, \delta_4, \delta_5, \dots\} \Rightarrow \omega(z) = \sum_{d \in \Delta} \frac{z^d}{d!}$$

- ▶ Erdős–Rényi graphs: Poisson degree distribution
- ▶ Scale-free network: $\mathbb{P}(k) \sim k^{-\gamma}$

Theorem. [D., Ravelomanana]

Phase transition is shifted with “weighted” allowed degrees

$$\omega(z) = \sum_{d \in \Delta} c_d z^d$$

Open question. Can we mimic the behaviour of scale-free network using a weighted EGF for set degree constraints?

Bianconi–Barabási approach

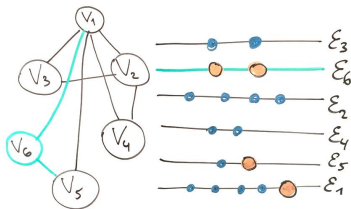
Bose–Einstein condensation in evolving networks

Bose gas

diluted gas at low temperature
temperature
energy
particle
number of energy levels
Bose–Einstein condensation

network evolution

evolving network
temperature
energy
half-edge
 \leq number of nodes
topological phase transition



In this model, the number of particles on the energy level ε follows the Bose statistics $n(\varepsilon) = \frac{1}{e^{\beta(\varepsilon-\mu)} - 1}$ which also represents the number of edges linking to nodes with energy ε .

Weighted partitions as BEC in quantum harmonic oscillator

[Vershik], [Bernstein, Fahrbach, Randall], [Bendkowski, Bodini, D.]

Coloured partitions \leftrightarrow d-dimensional quantum oscillator

Weighted partition

Sum of numbers

Number of colours

Row of Young table

Number of rows

Number of squares in the row

Partition limit shape

$$\binom{d+\lambda-1}{\lambda}$$

Random particle assembly

Total energy

Dimension (d)

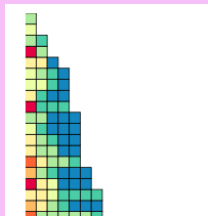
Particle

Number of particles

Energy of a particle (λ)

Bose–Einstein condensation

Number of particle states



Theorem. There exists a polynomial algorithm to generate random assemblies with given numbers of colours (n_1, n_2, \dots, n_d) .

Combinatorial learning

[D. '2020+]

Maximum likelihood approach. Data X_1, \dots, X_n is i.i.d, and follows \mathbb{P}_θ .

- ▶ Parameter estimation: find θ
- ▶ Hypothesis testing
- ▶ Nonlinear regression: $y = f(\mathbf{x}) + \varepsilon$
- ▶ Model selection, etc

Combinatorial learning. Data X_1, \dots, X_n are (projections of) parameters of combinatorial objects generated from Boltzmann distribution \mathbb{P}_z .

- ▶ Object size
- ▶ Number of leaves in a tree, height, etc
- ▶ Patterns in context-free grammars
- ▶ Basically, whatever is markable

Combinatorial learning

[D. '2020+]

Hidden parameter estimation. (Case study)

Objects are sampled from $\mathbb{P}_{z,u}$. We observe only the size parameter corresponding to z . Estimate (z, u) .

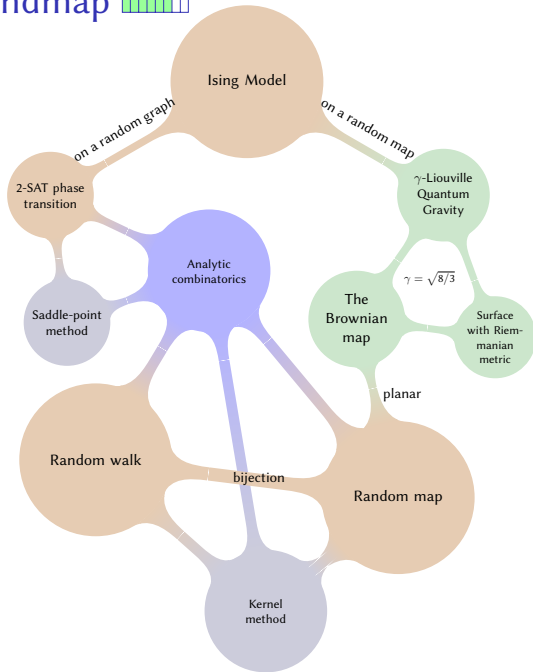
Lemma. In case of general dimension, this problem is #P-complete.

Proposition. Boltzmann relaxation has a polynomial solution, using the tuning algorithm from [Bendkowski, Bodini, D.]

Level III. Mathematical physics



Ising model mindmap



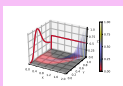
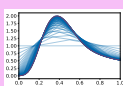
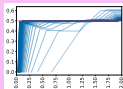
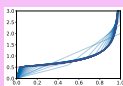
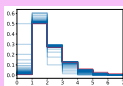
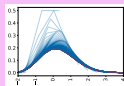
Connections with mathematical physics

- ▶ 2-SAT phase transition in terms of Ising model
- ▶ Maps, ribbon graphs and models of 2-d quantum gravity
- ▶ Weighted partitions and Bose–Einstein condensate

Parameters in random maps regardless of genus ▣▣▣▣▣▣▣▣

[Bodini, Courtiel, D., Hwang]

Statistics	Differential equation	Mean	Limit law
leaves	$L = v + (2 - u)zL + zL^2 + 2z^2\partial_z L + z(1 - v)\partial_v L$	1	Poisson(1)
root isthmus parts	$C = 1 + zC + v z C _{v=1} C + 2z^2\partial_z C$	2	Geometric($\frac{1}{2}$)
vertices	$X = v + zX + zX^2 + 2z^2\partial_z X$	$\log n$	$\mathcal{N}(\log n, \log n)$
loops	$Y = v + v z Y + v z Y _{v=1} Y + 2v z^2\partial_z Y + v^2 z(v w - 1)\partial_v Y$	$\frac{1}{2}n$	A new law*
root edges	$E = 1 + v z E + v z E _{v=1} E + 2v z^2\partial_z E$	$\frac{2}{3}n$	Beta($1, \frac{1}{2}$)
root degree	$D = 1 + v^2 z D + v z D _{v=1} D + 2v z^2\partial_z D - v^2(1 - v)z\partial_v D$	n	Uniform[0, 2]



Part III. Summary

Crossroad of tools and applications

	λ -calculus	networks	mathematical statistics	other
$F = \Phi(F, z)$	closed λ -terms			
$z \frac{F'(z)}{F(z)} = n$	software testing	Bianconi-Barabasi	combinatorial learning	tilings, RNA, queues
$P(F, F'_x, F'_y) = 0$	linear λ -terms			random maps
$\int_{\mathbb{R}} e^{-ix^3 - ixy} dx$	height profile	graphs with degree constraints		2-SAT Witten's conjecture

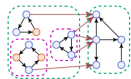
Paper



Shifting the phase transition threshold for random graphs using degree constraints



The birth of the contradictory component in random 2-SAT



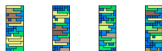
Symbolic method and directed graph enumeration



Statistical properties of lambda terms



Asymptotic distribution of parameters in random maps



Polynomial tuning of multiparametric combinatorial samplers

Tools

Combinatorial decomposition + saddle point

Inclusion-exclusion + saddle point

Combinatorial decomposition + Inclusion-exclusion

Infinite-dimensional Drmota–Lalley–Woods Theorem

Disconnecting transform + linearisation of PDE

Convex optimisation

Preprints and publications

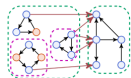
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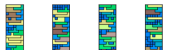
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- ▶ Peer-reviewed conference proceedings: 4 papers
- ▶ Accepted to a peer-reviewed journal: 1 paper ([70 pages, E-JC])

Preprints and publications

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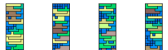
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Thank you.